

# An Approach to Quantify Reconfiguration Methods for PIM-SM Trees

## - A Network Complexity Perspective -

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# Introduction

## Motivation

- low predictability in network algorithms
- complete state space not transparent
- network algorithms also interact with underlying layers

## Complexity

- some complexity is needed for robust networks
- stability decreases at a certain level of complexity

# Aspects of Complexity

## Efficiency

- | classical measurement of effort for specified task

## Robustness

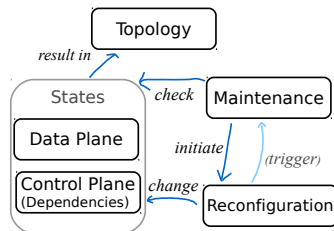
- | ability to persist despite changes

## Resilience

- | ability to adapt to changed conditions

# Control Loop

- behavior specification of a local node
- abstract view provides categorization of functions
- complexity arises from reconfiguration functions



# Related Work

## Stability of a Multicast Tree [1]

- quantifies the change of links in a multicast tree
- shows that the change follows a Poisson distribution
- formula to approximate the change of the tree depending on network size and number of receiver

## Percolation Thresholds in Network Decay [2, 3]

- describes characteristics of network stability
- high stability until a network decays completely
- exist in several network scenarios

# Current Work

## Goal

- estimate the range of change initiated by link failures
- estimate the ability of a network to adjust to failures

## PR 1 (work in progress)

- quantify effects of reconfiguration methods on distribution trees
- case study for PIM-SM [4]
  - join
  - reconfiguration at link failure
  - calculation based on expected state changes

# PIM-SM Protocol Overview

- IP layer multicast protocol
- independent of underlying routing protocol
- PIM routers maintain a variety of states
  - e.g. joins(\*,\*,RP(G)), prunes(S,G), lost\_assert(S,G), ...
- within two kinds of trees:
  - shared tree
  - source specific tree
- well known rendezvous point (RP) is root of the shared tree
- protocol consist of three main procedures (phases)



# PIM-SM Phases

## Phase One: RP Tree

- Receiver joins shared tree (`join(*,G)`)
- Sender packages are encapsulated sent to RP by DR

## Phase Two: Register Stop

- RP joins source specific tree (`join(S,G)`)
- When RP receives Packages via SST, send `register_stop(S,G)` to the sources DR

## Phase Three: Shortest Path Tree

- Receiver sends `join(S,G)` to source
- Receiver sends `prune(S,G,RPT)` to RP

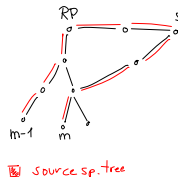
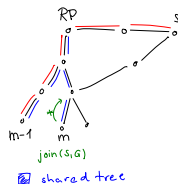
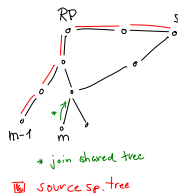
# Approach

- Given: approximation of the expected amount of new links in a multicast distribution tree  $E[\Delta_N(m)]$ 
  - if the  $m^{\text{th}}$  member joins (or leaves)
  - in a random network with  $N$  nodes
  - for big  $N$  and small  $m$

## Idea

Consider the reconfiguration operations, e.g. following a link failure, as a sum of several  $E[\Delta_N(m)]$ -like tree operations

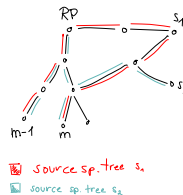
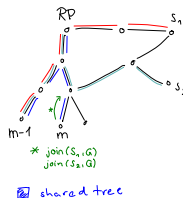
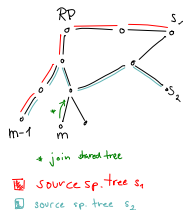
# Join Group with One Source



$m_{rp}$       number of members in shared tree  
 $m_s$       number of members in source specific tree

1 source:  $E[\Delta_N(m_{rp}, m_s)] \approx E[\Delta_N(m_{rp})] + E[\Delta_N(m_s)]$

# Join Group with Multiple Sources



$k$  sources:

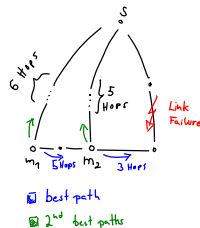
$$E[\Delta_N(m_{rp}, m_s)] \approx E[\Delta_N(m_{rp})] + k \cdot E[\Delta_N(m_s)] + \underbrace{k \cdot E[\Delta_N(m_{rp})]}_{\text{source specific prune}}$$

## Problem

- source specific prune within the join procedure cannot be approximated with  $E[\Delta_N(m_{rp})]$
- number of the receivers, which did not already prune a source, differs from  $m_{rp}$

# Effects of Link Failure

- link failure repaired by *multiple* rejoins
- How many nodes may be affected?
  - best case: 1
  - worst case: all nodes in the subtree



## Influencing Factors

- size and shape of the downstream subtree
- node degree of the upstream node

# Approach to Rejoin I

## Provisional Assumption

■ The existing distribution tree is a perfect k-ary tree.

Assuming a uniform link stability, the expected level of a link failure, can be calculated as

$$E(\ell) = \sum_{i=1}^h \frac{k^i}{N} \cdot i$$

( $\ell$ = level, N=number of links, k=downstream node degree, h=height of the tree)

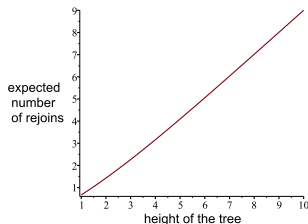
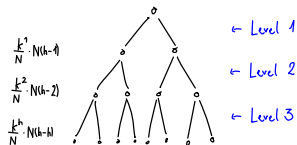
A first idea how to approach such a problem analytically!

# Approach to Rejoin II

Therefore the number of nodes which have to rejoin (avg. worst case) is:

$$E_{join}(h) = \sum_{i=1}^h \frac{k^i}{N(h)-1} \cdot N(h-i)$$

- $\frac{k^i}{N(h)-1}$  : probability to choose a link from level  $i$  from all links in the tree
- $N(x)$ : calculates the number of nodes from the number of levels  $x$



# Next Steps

- use a more realistic models to estimate the reach of reconfiguration methods (e.g. uniform recursive trees)
- add a time parameter to the calculations
- quantify the number of failing links per time, until the rejoin becomes ineffective
- check if the decay has a percolation threshold
- validate formulas against test results
  - e.g. from simulation or measurements



# Chances

- mathematical description for the effect of reconfiguration methods
- finding environments under which the algorithm will produce suboptimal results without bruteforce testing
- staying abstract enough to do similar calculation for other algorithms

# Risks

- interesting factors can be too difficult to model analytically
- too many expectancy values generalize the behavior to meaningless statements
- results do not validate in tests

# Summary

- reconfiguration methods are assumed to be strongly responsible for robustness
- framework of specified trees and  $E[\Delta_N(m)]$ -like operations enables to approximate affected nodes
- approach to quantify the robustness and effects of reconfigurations methods of PIM-SM
  - join
  - rejoin after link failure

# Future Work

- use of internet measurements for more detailed information on existing network structures
- check, how far analytical methods comply with measurements
- search for reasons why they would not comply
- consider not only multicast trees but tree creating algorithms in general (e.g. routing)

Thank you for your attention!

Questions?

# References I

- [1] P. Van Mieghem and M. Janic, "Stability of a Multicast Tree," in *INFOCOM 2002. Twenty-First Annual Joint Conference of the IEEE Computer and Communications Societies. Proceedings. IEEE*, vol. 2. IEEE, 2002, pp. 1099–1108.
- [2] R. Albert and A.-L. Barabási, "Statistical mechanics of complex networks," *Reviews of modern physics*, vol. 74, no. 1, p. 47, 2002.
- [3] S. V. Buldyrev, R. Parshani, G. Paul, H. E. Stanley, and S. Havlin, "Catastrophic cascade of failures in interdependent networks," *Nature*, vol. 464, no. 7291, pp. 1025–1028, 2010.
- [4] B. Fenner, M. Handley, H. Holbrook, and I. Kouvelas, "Protocol Independent Multicast - Sparse Mode (PIM-SM): Protocol Specification (Revised)," IETF, RFC 4601, August 2006.